

# Space-Time Chip Equalization for Maximum Diversity Space-Time Block Coded DS-CDMA Downlink Transmission

## Geert Leus

*Faculty of Electrical Engineering, Mathematics, and Computer Science, Delft University of Technology, Mekelweg 4, 2628CD Delft, The Netherlands  
Email: leus@cas.et.tudelft.nl*

## Frederik Petré

*Wireless Research, Interuniversity Micro-Electronics Center (IMEC), Kapeldreef 75, 3001 Leuven, Belgium  
Email: petre@imec.be*

## Marc Moonen

*Department of Electrical Engineering (ESAT), Katholieke Universiteit Leuven (K.U.Leuven), Kasteelpark Arenberg 10, 3001 Leuven, Belgium  
Email: moonen@esat.kuleuven.ac.be*

*Received 24 December 2002; Revised 4 August 2003*

In the downlink of DS-CDMA, frequency-selectivity destroys the orthogonality of the user signals and introduces multiuser interference (MUI). Space-time chip equalization is an efficient tool to restore the orthogonality of the user signals and suppress the MUI. Furthermore, multiple-input multiple-output (MIMO) communication techniques can result in a significant increase in capacity. This paper focuses on space-time block coding (STBC) techniques, and aims at combining STBC techniques with the original single-antenna DS-CDMA downlink scheme. This results into the so-called space-time block coded DS-CDMA downlink schemes, many of which have been presented in the past. We focus on a new scheme that enables both the maximum multiantenna diversity and the maximum multipath diversity. Although this maximum diversity can only be collected by maximum likelihood (ML) detection, we pursue suboptimal detection by means of space-time chip equalization, which lowers the computational complexity significantly. To design the space-time chip equalizers, we also propose efficient pilot-based methods. Simulation results show improved performance over the space-time RAKE receiver for the space-time block coded DS-CDMA downlink schemes that have been proposed for the UMTS and IS-2000 W-CDMA standards.

**Keywords and phrases:** downlink CDMA, space-time block coding, space-time chip equalization.

## 1. INTRODUCTION

Direct sequence code division multiple access (DS-CDMA) has emerged as the predominant multiple access technique for 3G cellular systems. In the downlink of DS-CDMA, orthogonal user signals are transmitted from the base station. All these signals are distorted by the same channel when propagating to the desired mobile station. Hence, when this channel is frequency-selective, the orthogonality of the user signals is destroyed and severe multiuser interference (MUI) is introduced. Space-time chip equalization can then restore the orthogonality of the user signals and suppress the MUI [1, 2, 3, 4].

Multiple-input multiple-output (MIMO) systems, on the other hand, have recently been shown to realize a significant

increase in capacity for rich scattering environments [5, 6, 7]. Both space division multiplexing (SDM) [8, 9] and space-time coding (STC) [10, 11, 12] are popular MIMO communication techniques. SDM techniques mainly aim at an increase in throughput by transmitting different data streams from the different transmit antennas. However, SDM typically requires as many receive as transmit antennas, which seriously impairs a cost-efficient implementation at the mobile station. STC techniques, on the other hand, mainly aim at an increase in performance by introducing spatial and temporal correlation in the transmitted data streams. As opposed to SDM, STC supports any number of receive antennas, and thus enables a cost-efficient implementation at the mobile station. In this perspective, space-time block coding (STBC) techniques, introduced in [11] for two transmit antennas and

later generalized in [12] for any number of transmit antennas, are particularly appealing because they facilitate maximum likelihood (ML) detection with simple linear processing. However, these STBC techniques have originally been developed for signaling over frequency-flat channels, and do not enable the maximum multiantenna and multipath diversity present in frequency-selective channels. Therefore, improved STBC techniques have recently been developed for signaling over frequency-selective channels [13, 14, 15]. The STBC technique proposed in [13] enables the maximum multiantenna diversity, and although it is presented as a technique that provides the maximum multipath diversity, it is not possible to prove it without any proper discussion on how to treat the edge effects at the beginning and the end of a burst. If the edge effects are handled by a cyclic prefix as in [14], maximum multipath diversity is not guaranteed. On the other hand, if the edge effects are handled by a zero post-fix as in [15], maximum multipath diversity is guaranteed.

Up till now, research on STBC techniques has mainly focused on single-user communication links. In this paper, we aim at combining STBC techniques with the original single-antenna DS-CDMA downlink scheme, resulting into so-called space-time block coded DS-CDMA downlink schemes. As an example, we mention the space-time block coded DS-CDMA downlink schemes that have been proposed for the UMTS and IS-2000 W-CDMA standards, both special cases of the so-called space-time spreading scheme presented in [16], which consists of a mixture of the original single-antenna DS-CDMA downlink scheme and the STBC technique of [12]. However, this scheme does not enable the maximum multiantenna and multipath diversity present in frequency-selective channels. A second example is the space-time block coded DS-CDMA downlink scheme presented in [17], which consists of the original single-antenna DS-CDMA downlink scheme followed by the STBC technique of [14]. However, this scheme only enables the maximum multiantenna diversity but not the maximum multipath diversity (due to the fact that maximum multipath diversity is not provided by the STBC technique of [14]). Therefore, in this paper, we consider the space-time block coded DS-CDMA downlink scheme that consists of the original single-antenna DS-CDMA downlink scheme followed by the STBC technique of [15]. This scheme enables both the maximum multiantenna diversity and the maximum multipath diversity (due to the fact that maximum multipath diversity is provided by the STBC technique of [15]). Although this maximum diversity can only be collected by ML detection, we pursue suboptimal detection by means of space-time chip equalization, which lowers the computational complexity significantly. Note that this suboptimal detection technique can also be applied to the STBC technique of [15] on its own, without combining it with the original single-antenna DS-CDMA downlink scheme.

Assuming there are  $J$  transmit antennas, the straightforward way to implement space-time chip equalization is to apply  $J$  space-time chip equalizers to recover the  $J$  transmitted space-time block coded multiuser chip sequences, then to apply space-time decoding to recover  $J$  subsequences of

the original multiuser chip sequence, and finally, to perform simple despreading. Since this comes down to an equalization problem with  $J$  sources, we need  $J + 1$  chip rate sampled outputs at each mobile station for a finite-length zero-forcing (ZF) solution to exist (i.e.,  $J + 1$  receive antennas if the antennas are sampled at chip rate). However, we will show that the space-time chip equalization and space-time decoding operations can be swapped, which allows us to first apply space-time decoding, then to apply  $J$  space-time chip equalizers to recover  $J$  subsequences of the original multiuser chip sequence, and finally, to perform simple despreading. Since this comes down to  $J$  equalization problems with only one source, we need only two chip rate sampled outputs at each mobile station for a finite-length ZF solution to exist (i.e., two receive antennas if the antennas are sampled at chip rate). To design the space-time chip equalizers, we finally propose efficient pilot-based methods.

In Section 2, we discuss the transceiver design of the proposed space-time block coded DS-CDMA system. We distinguish between the transmitter design, the channel model, and the receiver design, where the latter is based on space-time chip equalization. In Section 3, we then propose two pilot-based methods for practical space-time chip equalizer design. We show some simulation results in Section 4. In Section 5, we finally draw our conclusions.

### Notation

We use upper (lower) bold face letters to denote matrices (vectors). Superscripts  $*$ ,  $T$ , and  $H$  represent conjugate, transpose, and Hermitian, respectively. Further,  $\lfloor \cdot \rfloor$  represents the flooring operation, and  $\mathcal{E}\{\cdot\}$  represents the expectation operation. We denote the  $N \times N$  identity matrix as  $\mathbf{I}_N$  and the  $M \times N$  all-zero matrix as  $\mathbf{0}_{M \times N}$ . Next,  $[\mathbf{A}]_{m,n}$  denotes the entry at position  $(m, n)$  of the matrix  $\mathbf{A}$ . Finally,  $\text{diag}\{\mathbf{a}\}$  represents the diagonal matrix with the vector  $\mathbf{a}$  on the diagonal.

## 2. TRANSCIEVER DESIGN

We consider the downlink of a space-time block coded DS-CDMA system. We assume the base station is equipped with  $J$  transmit antennas, and the mobile station is equipped with  $M$  receive antennas. In the following, we discuss the transmitter design, the channel model, and the receiver design.

### 2.1. Transmitter design

At the base station, a space-time block coded DS-CDMA downlink scheme transforms  $\{s_u[k]\}_{u=1}^U$  and  $s_p[k]$ , where  $s_u[k]$  is the  $u$ th user's data symbol sequence and  $s_p[k]$  is the pilot symbol sequence, into  $J$  space-time block coded multiuser chip sequences  $\{u_j[n]\}_{j=1}^J$ .

We consider the space-time block coded DS-CDMA downlink scheme that consists of the original single-antenna CDMA downlink transmission scheme followed by the STBC technique of [15]. This scheme enables both the maximum multiantenna diversity and the maximum multipath diversity. For simplicity, we will focus on the case of  $J = 2$  transmit antennas. Extensions to more than two transmit antennas

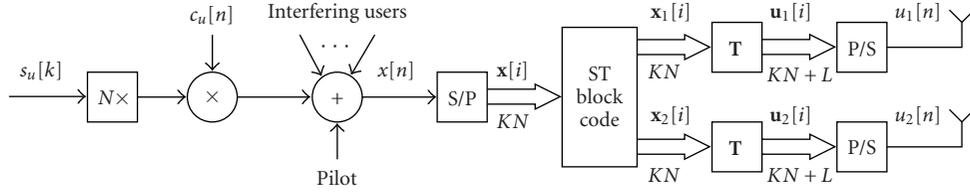


FIGURE 1: Proposed space-time block coded DS-CDMA downlink scheme.

( $J > 2$ ) are straightforward and can be developed following the design rules presented in [18].

Figure 1 depicts the proposed space-time block coded DS-CDMA downlink scheme ( $N \times$  repeats each sample  $N$  times, whereas “S/P” and “P/S” represent a serial-to-parallel and parallel-to-serial conversion, respectively). First, the original multiuser chip sequence  $x[n]$  is constructed:

$$x[n] := \sum_{u=1}^U s_u[\lfloor n/N \rfloor] c_u[n] + s_p[\lfloor n/N \rfloor] c_p[n], \quad (1)$$

where  $c_u[n]$  is the  $u$ th user’s code sequence and  $c_p[n]$  is the pilot code sequence. We assume that both  $c_u[n]$  and  $c_p[n]$  are normalized and consist of a multiplication of a user/pilot specific orthogonal Walsh-Hadamard spreading code of length  $N$  and a base-station specific long scrambling code. Note that the above pilot insertion technique is similar to the so-called common pilot channel (CPICH) [19] in forthcoming 3G systems. Second, the original multiuser chip sequence  $x[n]$  is serial-to-parallel converted into the  $1 \times KN$  multiuser chip block sequence  $\mathbf{x}[i]$ :

$$\mathbf{x}[i] := [x[iKN], \dots, x[(i+1)KN - 1]]. \quad (2)$$

Third, the multiuser chip block sequence  $\mathbf{x}[i]$  is transformed into the two  $1 \times KN$  block sequences  $\mathbf{x}_1[i]$  and  $\mathbf{x}_2[i]$ :

$$\begin{bmatrix} \mathbf{x}_1[2i] & \mathbf{x}_1[2i+1] \\ \mathbf{x}_2[2i] & \mathbf{x}_2[2i+1] \end{bmatrix} := \begin{bmatrix} \mathbf{x}[2i] & -\mathbf{x}^*[2i+1]\mathbf{P}_{KN} \\ \mathbf{x}[2i+1] & \mathbf{x}^*[2i]\mathbf{P}_{KN} \end{bmatrix}, \quad (3)$$

where  $\mathbf{P}_N$  is an  $N \times N$  permutation matrix that performs a reversal of the entries, that is,  $[\mathbf{P}_N]_{n,n'} = \delta[n + n' - N - 1]$ . Fourth, we add a zero postfix of length  $L$  to each block of the block sequence  $\mathbf{x}_j[i]$ , resulting into the  $1 \times (KN + L)$  block sequence  $\mathbf{u}_j[i]$ :  $\mathbf{u}_j[i] := \mathbf{x}_j[i]\mathbf{T}$ , where  $\mathbf{T}$  is the  $(KN) \times (KN + L)$  zero postfix insertion matrix:  $\mathbf{T} := [\mathbf{I}_{KN}, \mathbf{0}_{KN \times L}]$ . Finally, the block sequence  $\mathbf{u}_j[i]$  is parallel-to-serial converted into the space-time block coded multiuser chip sequence  $u_j[n]$ :

$$[u_j[i(KN + L)], \dots, u_j[(i+1)(KN + L) - 1]] := \mathbf{u}_j[i], \quad (4)$$

which is transmitted at the  $j$ th transmit antenna with rate  $1/T_c$  (the chip rate).

## 2.2. Channel model

Assuming the  $m$ th receive antenna is sampled at the chip rate, the received sequence at the  $m$ th receive antenna can be written as

$$y_m[n] = \sum_{j=1}^2 \sum_{l=0}^L h_{m,j}[l] u_j[n-l] + e_m[n], \quad (5)$$

where  $e_m[n]$  is the additive noise at the  $m$ th receive antenna and  $h_{m,j}[l]$  is the channel from the  $j$ th transmit antenna to the  $m$ th receive antenna, including transmit and receive filters. We assume that  $h_{m,j}[l]$  is FIR with order  $L_{j,m}$  and that  $L$  is a known upper bound on  $\max_{j,m} \{L_{j,m}\}$ . Note that  $L$  was also chosen as the zero postfix length in Section 2.1.

## 2.3. Receiver design

A first option is to serial-to-parallel convert the received sequence  $y_m[n]$  into the  $1 \times (KN + L)$  received block sequence  $\mathbf{y}_m[i]$ :

$$\mathbf{y}_m[i] := [y_m[i(KN + L)], \dots, y_m[(i+1)(KN + L) - 1]], \quad (6)$$

then to apply space-time decoding and Viterbi equalization as in [18], and finally, to perform simple despreading. This detection technique is overall ML, but leads to a very large computational complexity. That is why we pursue suboptimal detection by means of space-time chip equalization, which lowers the computational complexity significantly. Note that this suboptimal detection technique can also be applied to the STBC technique of [15] on its own, without combining it with the original single-antenna DS-CDMA downlink scheme.

We first introduce some new notation. Defining the  $M \times 1$  vector

$$\mathbf{y}[n] := [y_1[n], \dots, y_M[n]]^T, \quad (7)$$

we can write

$$\mathbf{y}[n] = \sum_{j=1}^2 \sum_{l=0}^L \mathbf{h}_j[l] u_j[n-l] + \mathbf{e}[n], \quad (8)$$

where  $\mathbf{e}[n]$  is similarly defined as  $\mathbf{y}[n]$ , and

$$\mathbf{h}_j[l] := [h_{1,j}[l], \dots, h_{M,j}[l]]^T. \quad (9)$$

Further, defining the  $(Q + 1)M \times KN$  matrix

$$\mathbf{Y}[i] := \begin{bmatrix} \mathbf{y}[i(KN + L)] & \cdots & \mathbf{y}[i(KN + L) + KN - 1] \\ \vdots & \vdots & \vdots \\ \mathbf{y}[i(KN + L) + Q] & \cdots & \mathbf{y}[i(KN + L) + KN - 1 + Q] \end{bmatrix}, \quad (10)$$

we can write

$$\mathbf{Y}[i] = \sum_{j=1}^2 \mathcal{H}_j \mathbf{U}_j[i] + \mathbf{E}[i], \quad (11)$$

where  $\mathbf{E}[i]$  is similarly defined as  $\mathbf{Y}[i]$ ,

$$\mathcal{H}_j := \begin{bmatrix} \mathbf{h}_j[L] & \cdots & \mathbf{h}_j[0] & \mathbf{0}_{M \times 1} & \cdots & \mathbf{0}_{M \times 1} \\ \mathbf{0}_{M \times 1} & \mathbf{h}_j[L] & \cdots & \mathbf{h}_j[0] & \cdots & \mathbf{0}_{M \times 1} \\ \vdots & \vdots & \ddots & \vdots & \ddots & \vdots \\ \mathbf{0}_{M \times 1} & \mathbf{0}_{M \times 1} & \cdots & \mathbf{h}_j[L] & \cdots & \mathbf{h}_j[0] \end{bmatrix},$$

$$\mathbf{U}_j[i] := \begin{bmatrix} u_j[i(KN + L) - L] & \cdots & u_j[i(KN + L) - L + KN - 1] \\ \vdots & \vdots & \vdots \\ u_j[i(KN + L) + Q] & \cdots & u_j[i(KN + L) + Q + KN - 1] \end{bmatrix}. \quad (12)$$

The parameter  $Q$  basically represents the order of the adopted space-time chip equalizer. This equalizer order  $Q$  is usually chosen to be close to the channel order  $L$ . For the sake of conciseness, we assume  $Q = L$ . However, the proposed results can easily be extended to other values of the equalizer order  $Q$ .

Choosing  $Q = L$ , it is clear from the zero postfix insertion that  $\mathbf{U}_j[i]$  can be expressed as

$$\mathbf{U}_j[i] = \mathcal{T}(\mathbf{x}_j[i]) := \begin{bmatrix} \mathbf{x}_j[i] \mathbf{J}_{KN}^{(-L)} \\ \vdots \\ \mathbf{x}_j[i] \mathbf{J}_{KN}^{(L)} \end{bmatrix}, \quad (13)$$

with  $\mathbf{J}_N^{(l)}$  the  $N \times N$  shift matrix with  $[\mathbf{J}_N^{(l)}]_{n,n'} = \delta[n - n' - l]$  (note that  $\mathbf{J}_N^{(0)} = \mathbf{I}_N$ ).

To proceed, the straightforward way is to apply two space-time chip equalizers on  $\mathbf{Y}[i]$  to recover  $\mathbf{x}_1[i]$  and  $\mathbf{x}_2[i]$ , then to apply space-time decoding to recover  $\mathbf{x}[2i]$  and  $\mathbf{x}[2i+1]$ , and finally, to perform simple despreading. Since this comes down to an equalization problem with two sources, we need three chip rate sampled receive antennas at each mobile station for a finite-length ZF solution to exist (for  $J > 2$  transmit antennas, we need  $J + 1$  chip rate sampled receive antennas at each mobile station). However, we will show that the space-time chip equalization and space-time decoding operations can be swapped, which allows us to first apply space-time decoding on  $\mathbf{Y}[2i]$  and  $\mathbf{Y}[2i+1]$ , then to apply two space-time chip equalizers to recover  $\mathbf{x}[2i]$  and  $\mathbf{x}[2i+1]$ , and finally, to perform simple despreading. Since this comes

down to two equalization problems with only one source, we need only two chip rate sampled receive antennas at each mobile station for a finite-length ZF solution to exist (even for  $J > 2$  transmit antennas, we need only two chip rate sampled receive antennas at each mobile station). The latter option clearly has more degrees of freedom to tackle the equalization problem, and therefore leads to a better performance. This option is explained in more detail next.

### 2.3.1. Space-time decoding

Using (11) and (13), we can write  $\mathbf{Y}[2i]$  and  $\mathbf{Y}[2i+1]$  as

$$\begin{aligned} \mathbf{Y}[2i] &= \mathcal{H}_1 \mathcal{T}(\mathbf{x}_1[2i]) + \mathcal{H}_2 \mathcal{T}(\mathbf{x}_2[2i]) + \mathbf{E}[2i], \\ \mathbf{Y}[2i+1] &= \mathcal{H}_1 \mathcal{T}(\mathbf{x}_1[2i+1]) + \mathcal{H}_2 \mathcal{T}(\mathbf{x}_2[2i+1]) \\ &\quad + \mathbf{E}[2i+1]. \end{aligned} \quad (14)$$

Since  $\mathbf{x}_1[2i+1] = -\mathbf{x}_2^*[2i] \mathbf{P}_{KN}$  (see (3)), we can derive from (13) that

$$\begin{aligned} \mathcal{T}(\mathbf{x}_1[2i+1]) &= \begin{bmatrix} \mathbf{x}_1[2i+1] \mathbf{J}_{KN}^{(-L)} \\ \vdots \\ \mathbf{x}_1[2i+1] \mathbf{J}_{KN}^{(L)} \end{bmatrix} \\ &= - \begin{bmatrix} \mathbf{x}_2^*[2i] \mathbf{P}_{KN} \mathbf{J}_{KN}^{(-L)} \\ \vdots \\ \mathbf{x}_2^*[2i] \mathbf{P}_{KN} \mathbf{J}_{KN}^{(L)} \end{bmatrix} \\ &= - \begin{bmatrix} \mathbf{x}_2^*[2i] \mathbf{J}_{KN}^{(L)} \\ \vdots \\ \mathbf{x}_2^*[2i] \mathbf{J}_{KN}^{(-L)} \end{bmatrix} \mathbf{P}_{KN} \\ &= -\mathbf{P}_{2L+1} \begin{bmatrix} \mathbf{x}_2^*[2i] \mathbf{J}_{KN}^{(-L)} \\ \mathbf{x}_2^*[2i] \mathbf{J}_{KN}^{(L)} \end{bmatrix} \mathbf{P}_{KN} \\ &= -\mathbf{P}_{2L+1} \mathcal{T}^*(\mathbf{x}_2[2i]) \mathbf{P}_{KN}. \end{aligned} \quad (15)$$

Similarly, since  $\mathbf{x}_2[2i+1] = \mathbf{x}_1^*[2i] \mathbf{P}_{KN}$  (see (3)), we can derive from (13) that

$$\mathcal{T}(\mathbf{x}_2[2i+1]) = \mathbf{P}_{2L+1} \mathcal{T}^*(\mathbf{x}_1[2i]) \mathbf{P}_{KN}. \quad (16)$$

Conjugating  $\mathbf{Y}[2i+1]$  and multiplying it to the right-hand side with  $\mathbf{P}_{KN}$ , we then arrive at

$$\begin{aligned} &\mathbf{Y}^*[2i+1] \mathbf{P}_{KN} \\ &= \mathcal{H}_1^* \mathcal{T}^*(\mathbf{x}_1[2i+1]) \mathbf{P}_{KN} + \mathcal{H}_2^* \mathcal{T}^*(\mathbf{x}_2[2i+1]) \mathbf{P}_{KN} \\ &\quad + \mathbf{E}^*[2i+1] \mathbf{P}_{KN} \\ &= -\mathcal{H}_1^* \mathbf{P}_{2L+1} \mathcal{T}(\mathbf{x}_2[2i]) + \mathcal{H}_2^* \mathbf{P}_{2L+1} \mathcal{T}(\mathbf{x}_1[2i]) \\ &\quad + \mathbf{E}^*[2i+1] \mathbf{P}_{KN}, \end{aligned} \quad (17)$$

where the second equality is due to (15) and (16). Stacking  $\mathbf{Y}[2i]$  and  $\mathbf{Y}^*[2i+1] \mathbf{P}_{KN}$ :

$$\tilde{\mathbf{Y}}[i] := \begin{bmatrix} \mathbf{Y}[2i] \\ \mathbf{Y}^*[2i+1] \mathbf{P}_{KN} \end{bmatrix}, \quad (18)$$

and using the fact that  $\mathbf{x}_1[2i] = \mathbf{x}[2i]$  and  $\mathbf{x}_2[2i] = \mathbf{x}[2i + 1]$  (see (3)), we finally obtain

$$\tilde{\mathbf{Y}}[i] = \mathcal{H}\tilde{\mathbf{X}}[i] + \tilde{\mathbf{E}}[i], \quad (19)$$

where  $\tilde{\mathbf{E}}[i]$  is similarly defined as  $\tilde{\mathbf{Y}}[i]$ ,

$$\mathcal{H} := \begin{bmatrix} \mathcal{H}_1 & \mathcal{H}_2 \\ \mathcal{H}_2^* \mathbf{P}_{2L+1} & -\mathcal{H}_1^* \mathbf{P}_{2L+1} \end{bmatrix}, \quad (20)$$

$$\tilde{\mathbf{X}}[i] := \begin{bmatrix} \mathcal{T}(\mathbf{x}[2i]) \\ \mathcal{T}(\mathbf{x}[2i + 1]) \end{bmatrix}.$$

### 2.3.2. Space-time chip equalization

We now apply two space-time chip equalizers on  $\tilde{\mathbf{Y}}[i]$ :  $\mathbf{f}_e$  and  $\mathbf{f}_o$ . The  $1 \times 2(L+1)M$  space-time chip equalizer  $\mathbf{f}_e$  is designed to extract the even multiuser chip block  $\mathbf{x}[2i]$ , whereas the  $1 \times 2(L+1)M$  space-time chip equalizer  $\mathbf{f}_o$  is designed to extract the odd multiuser chip block  $\mathbf{x}[2i + 1]$ :

$$\hat{\mathbf{x}}[2i] = \mathbf{f}_e \tilde{\mathbf{Y}}[i], \quad \hat{\mathbf{x}}[2i + 1] = \mathbf{f}_o \tilde{\mathbf{Y}}[i]. \quad (21)$$

Note that  $\mathbf{x}[2i]$  and  $\mathbf{x}[2i + 1]$  are two distinct rows of  $\tilde{\mathbf{X}}[i]$ .

A first possibility is to apply two ZF space-time chip equalizers, completely eliminating the interchip interference (ICI) at the expense of potentially excessive noise enhancement:

$$\mathbf{f}_e = \mathbf{i}_e (\mathcal{H}^H \mathbf{R}_e^{-1} \mathcal{H})^{-1} \mathcal{H}^H \mathbf{R}_e^{-1}, \quad (22)$$

$$\mathbf{f}_o = \mathbf{i}_o (\mathcal{H}^H \mathbf{R}_e^{-1} \mathcal{H})^{-1} \mathcal{H}^H \mathbf{R}_e^{-1},$$

where  $\mathbf{i}_e$  is a  $1 \times (4L+2)$  unit vector with a one in the  $(L+1)$ th position,  $\mathbf{i}_o$  is a  $1 \times (4L+2)$  unit vector with a one in the  $(3L+2)$ th position, and  $\mathbf{R}_e := 1/(KN) \mathcal{E}\{\tilde{\mathbf{E}}[i]\tilde{\mathbf{E}}^H[i]\}$ . A second possibility is to apply two minimum mean-squared error (MMSE) space-time chip equalizers, balancing ICI elimination with noise enhancement:

$$\mathbf{f}_e = \mathbf{i}_e (\mathcal{H}^H \mathbf{R}_e^{-1} \mathcal{H} + \mathbf{R}_x^{-1})^{-1} \mathcal{H}^H \mathbf{R}_e^{-1}, \quad (23)$$

$$\mathbf{f}_o = \mathbf{i}_o (\mathcal{H}^H \mathbf{R}_e^{-1} \mathcal{H} + \mathbf{R}_x^{-1})^{-1} \mathcal{H}^H \mathbf{R}_e^{-1},$$

where  $\mathbf{R}_x := 1/(KN) \mathcal{E}\{\tilde{\mathbf{X}}[i]\tilde{\mathbf{X}}^H[i]\}$ .

Assuming the additive noise sequences  $\{e_m[n]\}_{m=1}^M$  are mutually uncorrelated and white with variance  $\sigma_e^2$ , we can write  $\mathbf{R}_e = \sigma_e^2 \mathbf{I}_{2(L+1)M}$ . Furthermore, assuming the data symbol sequences  $\{s_u[n]\}_{u=1}^U$  are mutually uncorrelated and white with variance  $\sigma_s^2$ , the original multiuser chip sequence  $x[n]$  is white with variance  $\sigma_x^2 = \sigma_s^2 J/N$  (justified by the long scrambling code), and we can write  $\mathbf{R}_x = \sigma_x^2 \text{diag}\{\mathbf{r}_x, \mathbf{r}_x\} = \sigma_s^2 J/N \text{diag}\{\mathbf{r}_x, \mathbf{r}_x\}$ , where  $\mathbf{r}_x = [(KN-L)/(KN), \dots, (KN-1)/(KN), 1, (KN-1)/(KN), \dots, (KN-L)/(KN)]$ .

### 2.3.3. Despreading

We define the  $1 \times KU$  multiuser data symbol block  $\mathbf{s}[i]$  as

$$\mathbf{s}[i] := [\mathbf{s}_1[i], \dots, \mathbf{s}_U[i]], \quad (24)$$

where  $\mathbf{s}_u[i]$  is the  $u$ th user's  $1 \times K$  data symbol block given by

$$\mathbf{s}_u[i] := [s_u[iK], \dots, s_u[(i+1)K-1]]. \quad (25)$$

Note that the  $1 \times K$  pilot symbol block  $\mathbf{s}_p[i]$  is similarly defined as  $\mathbf{s}_u[i]$ . We further define the multiuser code matrix  $\mathbf{C}[i]$  as

$$\mathbf{C}[i] := [\mathbf{C}_1[i]^T, \dots, \mathbf{C}_U[i]^T]^T, \quad (26)$$

where  $\mathbf{C}_u[i]$  is the  $u$ th user's code matrix given by

$$\mathbf{C}_u[i] := \begin{bmatrix} \mathbf{c}_u[iK] & & \\ & \ddots & \\ & & \mathbf{c}_u[(i+1)K-1] \end{bmatrix}, \quad (27)$$

with  $\mathbf{c}_u[k] := [c_u[kN], \dots, c_u[(k+1)N-1]]$ . Note that the pilot code matrix  $\mathbf{C}_p[i]$  is similarly defined as  $\mathbf{C}_u[i]$ . It is then clear from (1) that the multiuser chip block  $\mathbf{x}[i]$  can be expressed as

$$\mathbf{x}[i] = \sum_{u=1}^U \mathbf{s}_u[i] \mathbf{C}_u[i] + \mathbf{s}_p[i] \mathbf{C}_p[i] \quad (28)$$

$$= \mathbf{s}[i] \mathbf{C}[i] + \mathbf{s}_p[i] \mathbf{C}_p[i].$$

Hence, by despreading the multiuser chip block  $\mathbf{x}[i]$  with the  $u$ th user's code matrix  $\mathbf{C}_u[i]$ , we obtain

$$\mathbf{s}_u[i] = \mathbf{x}[i] \mathbf{C}_u^H[i] \quad (29)$$

because  $\mathbf{C}_p[i] \mathbf{C}_u^H[i] = \mathbf{0}_{K \times K}$ ,  $\mathbf{C}_{u'}[i] \mathbf{C}_u^H[i] = \mathbf{0}_{K \times K}$  for  $u \neq u'$ , and  $\mathbf{C}_u[i] \mathbf{C}_u^H[i] = \mathbf{I}_K$ . Therefore, once  $\mathbf{x}[i]$  has been estimated, we can find an estimate for  $\mathbf{s}_u[i]$  by simple despreading:

$$\hat{\mathbf{s}}_u[i] = \hat{\mathbf{x}}[i] \mathbf{C}_u^H[i]. \quad (30)$$

Plugging (30) into (21), we thus obtain

$$\hat{\mathbf{s}}_u[2i] = \mathbf{f}_e \tilde{\mathbf{Y}}[i] \mathbf{C}_u^H[2i], \quad (31)$$

$$\hat{\mathbf{s}}_u[2i + 1] = \mathbf{f}_o \tilde{\mathbf{Y}}[i] \mathbf{C}_u^H[2i + 1].$$

From these equations, it is also clear that the order of equalization and despreading can be reversed. In other words, we can first despread  $\tilde{\mathbf{Y}}[i]$  with  $\mathbf{C}_u[2i]$  and  $\mathbf{C}_u[2i + 1]$ , and then perform space-time chip equalization on both results.

## 3. PRACTICAL SPACE-TIME CHIP EQUALIZER DESIGN

In this section, we focus on practical space-time chip equalizer design. In [20, 21], we have developed two pilot-based space-time chip equalizer design methods for the original single-antenna DS-CDMA downlink scheme: a *training-based* method and a *semiblind* method. In this section, these two methods are appropriately modified and applied to the

proposed space-time coded DS-CDMA downlink scheme. We consider a burst of  $2I$  data symbol blocks.

The goal of the *training-based* method is to compute the  $u$ th user's even and odd data symbol blocks  $\{\mathbf{s}_u[2i]\}_{i=1}^I$  and  $\{\mathbf{s}_u[2i+1]\}_{i=1}^I$  from  $\{\tilde{\mathbf{Y}}[i]\}_{i=1}^I$ , based on the even and odd pilot symbol blocks  $\{\mathbf{s}_p[2i]\}_{i=1}^I$  and  $\{\mathbf{s}_p[2i+1]\}_{i=1}^I$ , the even and odd pilot code matrices  $\{\mathbf{C}_p[2i]\}_{i=1}^I$  and  $\{\mathbf{C}_p[2i+1]\}_{i=1}^I$ , and the  $u$ th user's even and odd code matrices  $\{\mathbf{C}_u[2i]\}_{i=1}^I$  and  $\{\mathbf{C}_u[2i+1]\}_{i=1}^I$ .

The goal of the *semiblind* method is to compute the  $u$ th user's even and odd data symbol blocks  $\{\mathbf{s}_u[2i]\}_{i=1}^I$  and  $\{\mathbf{s}_u[2i+1]\}_{i=1}^I$  from  $\{\tilde{\mathbf{Y}}[i]\}_{i=1}^I$ , based on the even and odd pilot symbol blocks  $\{\mathbf{s}_p[2i]\}_{i=1}^I$  and  $\{\mathbf{s}_p[2i+1]\}_{i=1}^I$ , the even and odd pilot code matrices  $\{\mathbf{C}_p[2i]\}_{i=1}^I$  and  $\{\mathbf{C}_p[2i+1]\}_{i=1}^I$ , and the even and odd multiuser code matrices  $\{\mathbf{C}[2i]\}_{i=1}^I$  and  $\{\mathbf{C}[2i+1]\}_{i=1}^I$ . Note that the semiblind method requires the knowledge of the active codes. This knowledge can be obtained by means of a limited feedback from the base station to the mobile station (only the indices of the active codes have to be fed back). However, this knowledge can also be obtained by first adopting the training-based method to design a space-time chip equalizer, and then comparing for each code the energy obtained after equalization and despreading with some threshold in order to decide whether this code is active or not.

For the sake of conciseness, we will only focus on block implementations. These block implementations might look rather complex, but they form the basis for practical low-complexity adaptive implementations, which can be derived in a similar fashion as done in [20, 21].

For the sake of simplicity, we make the following assumptions:

- (A1) the matrix  $\mathcal{H}$  has full column rank  $4L + 2$ ;
- (A2) the matrices  $\tilde{\mathbf{X}}[2i]$  and  $\tilde{\mathbf{X}}[2i + 1]$  have full row rank  $4L + 2$  for all  $i \in \{1, \dots, I\}$ .

The first assumption requires that  $2(L + 1)(M - 1) \geq 2L$ , which means we need only  $M \geq 2$  receive antennas at each mobile station (even for  $J > 2$  transmit antennas, we need only  $M \geq 2$  receive antennas at each mobile station). The second assumption requires that  $4L + 2 \leq KN$ . Note that these assumptions are not really necessary for the proposed methods to work. The only true requirement is that  $\mathbf{x}[2i]$  and  $\mathbf{x}[2i + 1]$  belong to the row space of  $\tilde{\mathbf{Y}}[i]$  for all  $i \in \{1, \dots, I\}$ . Assumptions (A1) and (A2) are sufficient but not necessary conditions for this. However, they considerably simplify the analysis.

Assume no noise is present. Because of assumption (A1), the row space of  $\tilde{\mathbf{Y}}[i]$  equals the row space of  $\tilde{\mathbf{X}}[i]$ . Hence, there exist two  $1 \times 2(L + 1)M$  space-time chip equalizers  $\hat{\mathbf{f}}_e$  and  $\hat{\mathbf{f}}_o$ , for which

$$\begin{aligned} \hat{\mathbf{f}}_e \tilde{\mathbf{Y}}[i] - \mathbf{x}[2i] &= \mathbf{0}_{1 \times KN}, \\ \hat{\mathbf{f}}_o \tilde{\mathbf{Y}}[i] - \mathbf{x}[2i + 1] &= \mathbf{0}_{1 \times KN}. \end{aligned} \quad (32)$$

Because of assumption (A2), these two space-time chip

equalizers  $\hat{\mathbf{f}}_e$  and  $\hat{\mathbf{f}}_o$  are ZF. By using (28), we then obtain

$$\begin{aligned} \hat{\mathbf{f}}_e \tilde{\mathbf{Y}}[i] - \mathbf{s}[2i] \mathbf{C}[2i] - \mathbf{s}_p[2i] \mathbf{C}_p[2i] &= \mathbf{0}_{1 \times KN}, \\ \hat{\mathbf{f}}_o \tilde{\mathbf{Y}}[i] - \mathbf{s}[2i + 1] \mathbf{C}[2i + 1] - \mathbf{s}_p[2i + 1] \mathbf{C}_p[2i + 1] &= \mathbf{0}_{1 \times KN}. \end{aligned} \quad (33)$$

### 3.1. Training-based method

By despreading (33) with the even and odd pilot code matrices  $\mathbf{C}_p[2i]$  and  $\mathbf{C}_p[2i + 1]$ , we obtain

$$\begin{aligned} \hat{\mathbf{f}}_e \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i] - \mathbf{s}_p[2i] &= \mathbf{0}_{1 \times K}, \\ \hat{\mathbf{f}}_o \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i + 1] - \mathbf{s}_p[2i + 1] &= \mathbf{0}_{1 \times K} \end{aligned} \quad (34)$$

because  $\mathbf{C}[i] \mathbf{C}_p^H[i] = \mathbf{0}_{K \times K}$  and  $\mathbf{C}_p[i] \mathbf{C}_p^H[i] = \mathbf{I}_K$ . The training-based method solves (34) for  $\hat{\mathbf{f}}_e$  and  $\hat{\mathbf{f}}_o$  for all  $i \in \{1, \dots, I\}$ . In the noisy case, this leads to the following least squares (LS) problems:

$$\begin{aligned} \min_{\hat{\mathbf{f}}_e} & \left\{ \sum_{i=1}^I \|\hat{\mathbf{f}}_e \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i] - \mathbf{s}_p[2i]\|^2 \right\}, \\ \min_{\hat{\mathbf{f}}_o} & \left\{ \sum_{i=1}^I \|\hat{\mathbf{f}}_o \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i + 1] - \mathbf{s}_p[2i + 1]\|^2 \right\}, \end{aligned} \quad (35)$$

which can be interpreted as follows. The space-time decoded output matrix  $\tilde{\mathbf{Y}}[i]$  is first equalized with the even and odd space-time chip equalizers  $\hat{\mathbf{f}}_e$  and  $\hat{\mathbf{f}}_o$ , and then despread with the even and odd pilot code matrices  $\mathbf{C}_p[2i]$  and  $\mathbf{C}_p[2i + 1]$ . The resulting even and odd vectors  $\hat{\mathbf{f}}_e \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i]$  and  $\hat{\mathbf{f}}_o \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i + 1]$  should then be as close as possible in an LS sense to the even and odd pilot symbol blocks  $\mathbf{s}_p[2i]$  and  $\mathbf{s}_p[2i + 1]$  for all  $i \in \{1, \dots, I\}$ . The solutions of (35) can be written as

$$\begin{aligned} \hat{\mathbf{f}}_e &= \left( \sum_{i=1}^I \mathbf{s}_p[2i] \mathbf{C}_p[2i] \tilde{\mathbf{Y}}^H[i] \right) \\ & \times \left( \sum_{i=1}^I \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i] \mathbf{C}_p[2i] \tilde{\mathbf{Y}}^H[i] \right)^{-1}, \\ \hat{\mathbf{f}}_o &= \left( \sum_{i=1}^I \mathbf{s}_p[2i + 1] \mathbf{C}_p[2i + 1] \tilde{\mathbf{Y}}^H[i] \right) \\ & \times \left( \sum_{i=1}^I \tilde{\mathbf{Y}}[i] \mathbf{C}_p^H[2i + 1] \mathbf{C}_p[2i + 1] \tilde{\mathbf{Y}}^H[i] \right)^{-1}. \end{aligned} \quad (36)$$

The obtained space-time chip equalizers  $\hat{\mathbf{f}}_e$  and  $\hat{\mathbf{f}}_o$  are subsequently used to estimate the  $u$ th user's even and odd data symbol blocks  $\hat{\mathbf{s}}_u[2i]$  and  $\hat{\mathbf{s}}_u[2i + 1]$  for all  $i \in \{1, \dots, I\}$ :

$$\begin{aligned} \hat{\mathbf{s}}_u[2i] &= \hat{\mathbf{f}}_e \tilde{\mathbf{Y}}[i] \mathbf{C}_u^H[2i], \\ \hat{\mathbf{s}}_u[2i + 1] &= \hat{\mathbf{f}}_o \tilde{\mathbf{Y}}[i] \mathbf{C}_u^H[2i + 1]. \end{aligned} \quad (37)$$

These soft estimates are fed into a decision device that determines the nearest constellation point.

### 3.2. Semiblind method

The semiblind method directly solves (33) for  $(\mathbf{f}_e, \mathbf{s}[2i])$  and  $(\mathbf{f}_o, \mathbf{s}[2i+1])$  for all  $i \in \{1, \dots, I\}$ . In the noisy case, this leads to the following LS problems:

$$\begin{aligned} \min_{(\mathbf{f}_e, \{\mathbf{s}[2i]\}_{i=1}^I)} & \left\{ \sum_{i=1}^I \|\mathbf{f}_e \bar{\mathbf{Y}}[i] - \mathbf{s}[2i] \mathbf{C}[2i] - \mathbf{s}_p[2i] \mathbf{C}_p[2i]\|^2 \right\}, \\ \min_{(\mathbf{f}_o, \{\mathbf{s}[2i+1]\}_{i=1}^I)} & \left\{ \sum_{i=1}^I \|\mathbf{f}_o \bar{\mathbf{Y}}[i] - \mathbf{s}[2i+1] \mathbf{C}[2i+1] \right. \\ & \left. - \mathbf{s}_p[2i+1] \mathbf{C}_p[2i+1]\|^2 \right\}. \end{aligned} \quad (38)$$

Since we are interested in  $\mathbf{f}_e$  and  $\mathbf{f}_o$ , we can first solve (38) for  $\hat{\mathbf{s}}[2i]$  and  $\hat{\mathbf{s}}[2i+1]$  for all  $i \in \{1, \dots, I\}$ , which results into

$$\begin{aligned} \hat{\mathbf{s}}[2i] &= \mathbf{f}_e \bar{\mathbf{Y}}[i] \mathbf{C}^H[2i], \\ \hat{\mathbf{s}}[2i+1] &= \mathbf{f}_o \bar{\mathbf{Y}}[i] \mathbf{C}^H[2i+1] \end{aligned} \quad (39)$$

because  $\mathbf{C}[i] \mathbf{C}_p^H[i] = \mathbf{0}_{K \times K}$  and  $\mathbf{C}_p[i] \mathbf{C}_p^H[i] = \mathbf{I}_K$ . Substituting  $\hat{\mathbf{s}}[2i]$  and  $\hat{\mathbf{s}}[2i+1]$  in (38) leads to the following LS problems:

$$\begin{aligned} \min_{\mathbf{f}_e} & \left\{ \sum_{i=1}^I \|\mathbf{f}_e \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i] \mathbf{C}[2i]) - \mathbf{s}_p[2i] \mathbf{C}_p[2i]\|^2 \right\}, \\ \min_{\mathbf{f}_o} & \left\{ \sum_{i=1}^I \|\mathbf{f}_o \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i+1] \mathbf{C}[2i+1]) \right. \\ & \left. - \mathbf{s}_p[2i+1] \mathbf{C}_p[2i+1]\|^2 \right\}, \end{aligned} \quad (40)$$

which can be interpreted as follows. The space-time decoded output matrix  $\bar{\mathbf{Y}}[i]$  is first equalized with the even and odd space-time chip equalizers  $\mathbf{f}_e$  and  $\mathbf{f}_o$  and then projected on the orthogonal complement of the subspace spanned by the even and odd multiuser code matrices  $\mathbf{C}[2i]$  and  $\mathbf{C}[2i+1]$ . The resulting even and odd vectors  $\mathbf{f}_e \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i] \mathbf{C}[2i])$  and  $\mathbf{f}_o \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i+1] \mathbf{C}[2i+1])$  should then be as close as possible in an LS sense to the even and odd pilot chip blocks  $\mathbf{s}_p[2i] \mathbf{C}_p[2i]$  and  $\mathbf{s}_p[2i+1] \mathbf{C}_p[2i+1]$  for all  $i \in \{1, \dots, I\}$ . The solutions of (40) can be written as

$$\begin{aligned} \hat{\mathbf{f}}_e &= \left( \sum_{i=1}^I \mathbf{s}_p[2i] \mathbf{C}_p[2i] \bar{\mathbf{Y}}^H[i] \right) \\ & \times \left( \sum_{i=1}^I \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i] \mathbf{C}[2i]) \bar{\mathbf{Y}}^H[i] \right)^{-1}, \\ \hat{\mathbf{f}}_o &= \left( \sum_{i=1}^I \mathbf{s}_p[2i+1] \mathbf{C}_p[2i+1] \bar{\mathbf{Y}}^H[i] \right) \\ & \times \left( \sum_{i=1}^I \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i+1] \mathbf{C}[2i+1]) \bar{\mathbf{Y}}^H[i] \right)^{-1}. \end{aligned} \quad (41)$$

The obtained space-time chip equalizers  $\hat{\mathbf{f}}_e$  and  $\hat{\mathbf{f}}_o$  are subsequently used to estimate the  $u$ th user's even and odd data symbol blocks  $\mathbf{s}_u[2i]$  and  $\mathbf{s}_u[2i+1]$  for all  $i \in \{1, \dots, I\}$ :

$$\begin{aligned} \hat{\mathbf{s}}_u[2i] &= \hat{\mathbf{f}}_e \bar{\mathbf{Y}}[i] \mathbf{C}_u^H[2i], \\ \hat{\mathbf{s}}_u[2i+1] &= \hat{\mathbf{f}}_o \bar{\mathbf{Y}}[i] \mathbf{C}_u^H[2i+1]. \end{aligned} \quad (42)$$

These soft estimates are fed into a decision device that determines the nearest constellation point.

With some algebraic manipulations, it is easy to prove that (40) is equivalent to

$$\begin{aligned} \min_{\mathbf{f}_e} & \left\{ \sum_{i=1}^I \|\mathbf{f}_e \bar{\mathbf{Y}}[i] \mathbf{C}_p^H[2i] - \mathbf{s}_p[2i]\|^2 \right. \\ & \left. + \|\mathbf{f}_e \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i] \mathbf{C}[2i] - \mathbf{C}_p^H[2i] \mathbf{C}_p[2i])\|^2 \right\}, \\ \min_{\mathbf{f}_o} & \left\{ \sum_{i=1}^I \|\mathbf{f}_o \bar{\mathbf{Y}}[i] \mathbf{C}_p^H[2i+1] - \mathbf{s}_p[2i+1]\|^2 \right. \\ & \left. + \|\mathbf{f}_o \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i+1] \mathbf{C}[2i+1] \right. \\ & \left. - \mathbf{C}_p^H[2i+1] \mathbf{C}_p[2i+1])\|^2 \right\}. \end{aligned} \quad (43)$$

This shows that (40) naturally decouples into a training-based part and a blind part (hence the name *semiblind*). The training-based part corresponds to (35). The blind part can be interpreted as follows. The space-time decoded output matrix  $\bar{\mathbf{Y}}[i]$  is first equalized with the even and odd space-time chip equalizers  $\mathbf{f}_e$  and  $\mathbf{f}_o$  and then projected on the orthogonal complement of the subspace spanned by the even and odd multiuser code matrices  $\mathbf{C}[2i]$  and  $\mathbf{C}[2i+1]$  and the even and odd pilot code matrices  $\mathbf{C}_p[2i]$  and  $\mathbf{C}_p[2i+1]$ . The resulting even and odd vectors  $\mathbf{f}_e \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i] \mathbf{C}[2i] - \mathbf{C}_p^H[2i] \mathbf{C}_p[2i])$  and  $\mathbf{f}_o \bar{\mathbf{Y}}[i] (\mathbf{I}_{KN} - \mathbf{C}^H[2i+1] \mathbf{C}[2i+1] - \mathbf{C}_p^H[2i+1] \mathbf{C}_p[2i+1])$  should then be as small as possible in an LS sense for all  $i \in \{1, \dots, I\}$ . Note that when the user load increases, the orthogonal complement of the subspace spanned by the even and odd multiuser code matrices  $\mathbf{C}[2i]$  and  $\mathbf{C}[2i+1]$  and the even and odd pilot code matrices  $\mathbf{C}_p[2i]$  and  $\mathbf{C}_p[2i+1]$  decreases in dimension. As a result, the information that the blind part contributes to the training-based part diminishes, and the semiblind method converges to the training-based method. In the extreme case when the system is fully loaded, that is,  $N = U - 1$ , the orthogonal complement of the subspace spanned by the even and odd multiuser code matrices  $\mathbf{C}[2i]$  and  $\mathbf{C}[2i+1]$  and the even and odd pilot code matrices  $\mathbf{C}_p[2i]$  and  $\mathbf{C}_p[2i+1]$  is empty, that is,  $\mathbf{I}_{KN} - \mathbf{C}^H[2i] \mathbf{C}[2i] - \mathbf{C}_p^H[2i] \mathbf{C}_p[2i] = \mathbf{0}_{KN \times KN}$  and  $\mathbf{I}_{KN} - \mathbf{C}^H[2i+1] \mathbf{C}[2i+1] - \mathbf{C}_p^H[2i+1] \mathbf{C}_p[2i+1] = \mathbf{0}_{KN \times KN}$ . Hence, the blind part does not contribute any additional information to the training-based part, and the semiblind method reduces to the training based method, that is, (43) reduces to (35).

#### 4. SIMULATION RESULTS

In this section, we compare the proposed space-time chip equalizer for the proposed space-time coded downlink CDMA transmission scheme with the space-time RAKE receiver for the space-time spreading scheme, which encompasses the space-time coded downlink CDMA transmission schemes that have been proposed for the UMTS and IS-2000 W-CDMA standards [16]. We do not consider channel codes when comparing the above transceivers. Otherwise, it will not be very clear whether a performance gain is due to the transceiver or the channel code. Moreover, the influence of channel codes on performance has been studied extensively in literature. In W-CDMA, the target coded BER typically is  $10^{-6}$ , which boils down to an uncoded BER of  $10^{-2}$  with a convolutional code of rate 1/2, constraint length 7, and soft decision Viterbi [22]. Therefore, we compare the different transceivers at an uncoded BER of  $10^{-2}$  in the sequel.

We consider a downlink CDMA system with a spreading factor of  $N = 32$ ,  $J = 2$  transmit antennas at the base station, and  $M = 2$  receive antennas at each mobile station. We assume that all channels are independent. We further assume that each channel  $h_{j,m}[n]$  is FIR with order  $L_{j,m} = 3$  and has independent Rayleigh fading channel taps of equal variance  $\sigma_h^2$ . Note that the bandwidth efficiency of the proposed space-time coded downlink CDMA transmission scheme is  $\epsilon_1 = KU/(KN + L)$ , whereas the bandwidth efficiency of the space-time spreading scheme is  $\epsilon_2 = U/N$ . Hence, in order to make a fair comparison between the two systems, their spectral efficiencies should be comparable. We therefore take  $K = 5$  and  $L = 3$  for the proposed space-time coded downlink CDMA transmission scheme, which results into  $\epsilon_1/\epsilon_2 \approx 0.98$ . We assume QPSK modulated data symbols, and define the signal-to-noise ratio (SNR) as the received bit energy over the noise power:

$$\begin{aligned} \text{SNR} &= \frac{\sigma_s^2/2 \sum_{j=1}^2 \sum_{l=0}^L \mathcal{E} \{ \|\mathbf{h}_j[l]\|^2 \}}{\sigma_e^2} \\ &= \frac{2(L+1)\sigma_s^2\sigma_h^2}{\sigma_e^2}. \end{aligned} \quad (44)$$

Two test cases are investigated.

##### Test case 1

We first assume that the pilot enables us to obtain perfect channel knowledge at the receiver. We then compare the proposed MMSE space-time chip equalizer for the proposed space-time coded downlink CDMA transmission scheme with the MMSE space-time RAKE receiver for the space-time spreading scheme (see [23, 24]), which is different from the matched space-time RAKE receiver for the space-time spreading scheme (see [16]) because it uses an MMSE filter instead of a matched filter to combine the finger outputs. It has been shown in [23, 24] that for the space-time spreading scheme, the MMSE space-time RAKE receiver significantly outperforms the matched space-time RAKE receiver. Figures 2, 3, and 4 compare the performance of the two transceivers

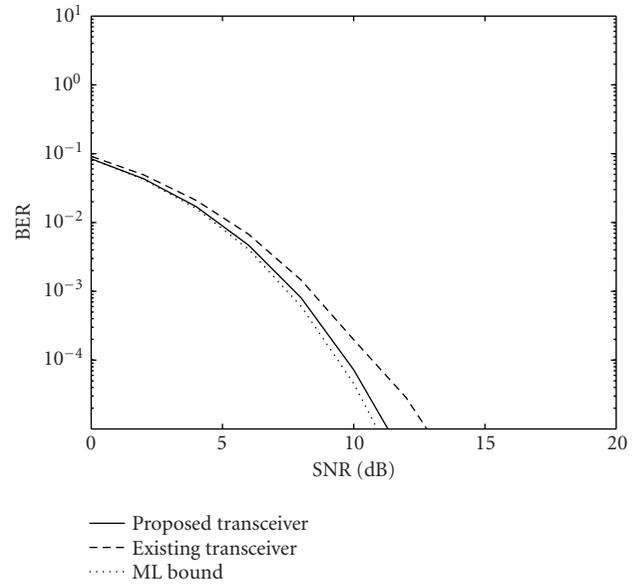


FIGURE 2: Performance comparison for  $U = 1$ .

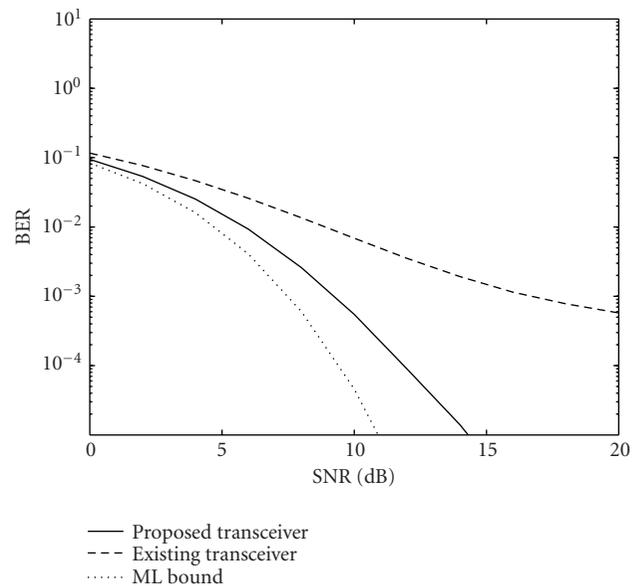
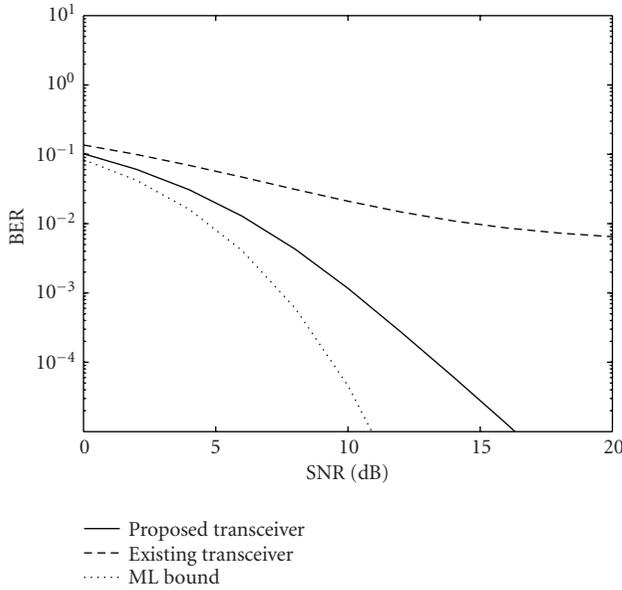
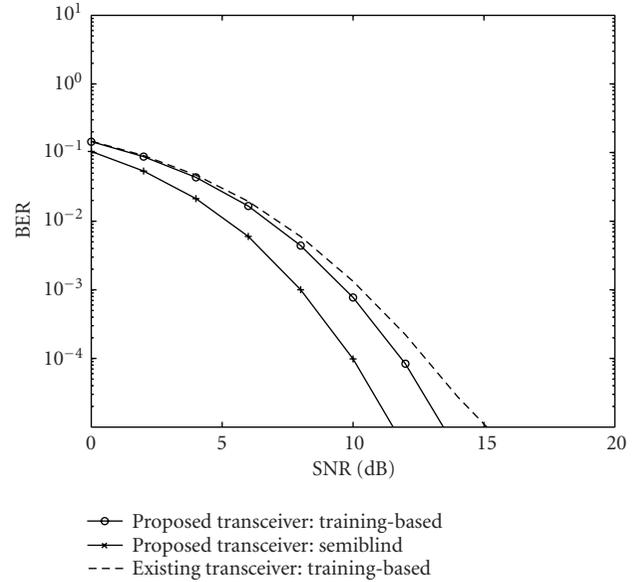


FIGURE 3: Performance comparison for  $U = 15$ .

for  $U = 1$ ,  $U = 15$ , and  $U = 31$  users, respectively. The performance results are averaged over 1000 random channel realizations, where for each channel realization, we consider 10 random data and noise realizations corresponding to  $I = 10$  (100 data symbols per user). Also shown is the theoretical performance of  $\sum_{j,m} (L_{j,m} + 1) = 16$ -fold diversity over Rayleigh fading channels [22].

First of all, we see that the proposed transceiver comes close to extracting the maximum diversity at low-to-medium

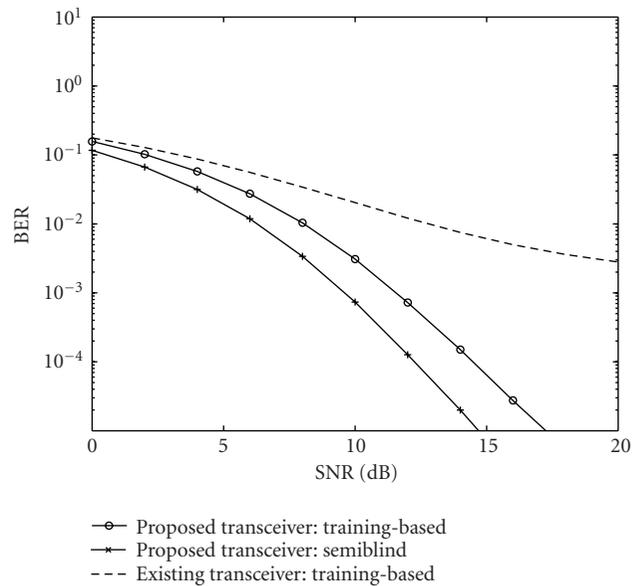
FIGURE 4: Performance comparison for  $U = 31$ .FIGURE 5: Performance of pilot-based methods for  $U = 1$ .

user loads. More specifically, at a BER of  $10^{-2}$ , the proposed transceiver incurs a 0.1, 1, and 1.8 dB loss compared to the theoretical ML bound for  $U = 1$ ,  $U = 15$ , and  $U = 31$  users, respectively. The existing transceiver, on the other hand, performs poorly at medium-to-high user loads. At a BER of  $10^{-2}$ , it incurs a 0.5, 3, and 8.2 dB performance loss compared to the proposed transceiver for  $U = 1$ ,  $U = 15$ , and  $U = 31$  users, respectively. The existing transceiver is not capable of completely suppressing the MUI at high SNR. This results into a flooring of the BER at high SNR. Note that the flooring level increases with the number of users  $U$ .

### Test case 2

We now investigate the performance of the pilot-based methods. Note that for the space-time spreading scheme, it is easy to derive a training-based method to estimate the combining filter of the space-time RAKE receiver based on the knowledge of the pilot. The performance results are again averaged over 1000 random channel realizations, where for each channel realization, we consider 10 random data and noise realizations corresponding to  $I = 10$  (100 data symbols per user). Figures 5, 6, and 7 compare the performance of the different methods for  $U = 1$ ,  $U = 15$ , and  $U = 31$  users, respectively.

First of all, we observe that the difference between the training-based method and the semiblind method for the proposed transceiver decreases with an increasing user load, as indicated in Section 3.2. Next, we observe that the training-based method for the existing transceiver performs much worse than the training-based and semiblind methods for the proposed transceiver at medium-to-high user loads. Finally, note that for the proposed transceiver, the MMSE performance discussed in test case 1 can be viewed

FIGURE 6: Performance of pilot-based methods for  $U = 15$ .

as the convergence point of the training-based and semiblind methods as  $I$  goes to infinity. Comparing the figures of test case 2 with the figures of test case 1, we observe that for  $I = 10$ , the training-based method is still far from the MMSE performance, whereas the semiblind method is already very close to the MMSE performance. Hence, as  $I$  increases, the semiblind method converges faster to the MMSE performance than the training-based method.

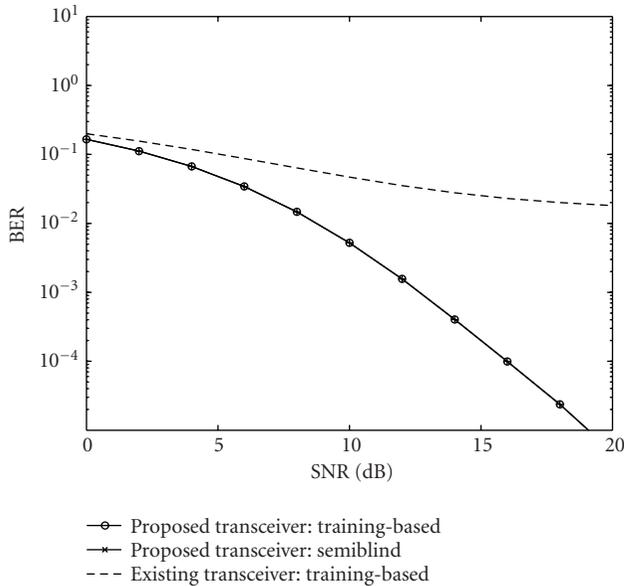


FIGURE 7: Performance of pilot-based methods for  $U = 31$ .

## 5. CONCLUSIONS

We have aimed at combining STBC techniques with the original single-antenna DS-CDMA downlink scheme, resulting into the so-called space-time block coded DS-CDMA downlink schemes. Many space-time block coded DS-CDMA downlink transmission schemes can be considered. We have focussed on a new scheme that enables both the maximum multiantenna diversity and the maximum multipath diversity. Although this maximum diversity can only be collected by ML detection, we have pursued suboptimal detection by means of space-time chip equalization, which lowers the computational complexity significantly. To design the space-time chip equalizers, we have also proposed efficient pilot-based methods. Simulation results have shown improved performance over the space-time RAKE receiver for the space-time block coded DS-CDMA downlink schemes that have been proposed for the UMTS and IS-2000 W-CDMA standards.

## ACKNOWLEDGMENTS

This research work was carried out in the frame of the Belgian State's Interuniversity Poles of Attraction Programme (2002–2007): IAP P5/22 (“Dynamical Systems and Control: Computation, Identification, and Modelling”) and P5/11 (“Mobile Multimedia Communication Systems and Networks”); the Concerted Research Action GOA-MEFISTO-666 (Mathematical Engineering for Information and Communication Systems Technology) of the Flemish Government; and Research Project FWO no. G.0196.02 (“Design of Efficient Communication Techniques for Wireless Time-Dispersive Multiuser MIMO Systems”). Part of this work ap-

peared in the proceedings of the International Conference on Communications (ICC), New York city, NY, April-May 2002. During this research work, Geert Leus was a Postdoctoral Fellow of the Fund for Scientific Research - Flanders (FWO - Vlaanderen), and Frederik Petré was a Research Assistant of the Institute for the Promotion of Innovation by Science and Technology in Flanders (IWT).

## REFERENCES

- [1] A. Klein, “Data detection algorithms specially designed for the downlink of CDMA mobile radio systems,” in *Proc. IEEE Vehicular Technology Conference*, vol. 1, pp. 203–207, Phoenix, Ariz, USA, May 1997.
- [2] I. Ghauri and D. Slock, “Linear receivers for the DS-CDMA downlink exploiting orthogonality of spreading sequences,” in *Proc. 32nd Asilomar Conf. on Signals, Systems, and Computers*, vol. 1, pp. 650–654, Pacific Grove, Calif, USA, November 1998.
- [3] T. P. Krauss, W. J. Hillery, and M. D. Zoltowski, “Downlink specific linear equalization for frequency selective CDMA cellular systems,” *Journal of VLSI Signal Processing*, vol. 30, no. 3, pp. 143–161, 2002.
- [4] C. D. Frank, E. Visotsky, and U. Madhow, “Adaptive interference suppression for the downlink of a direct sequence CDMA system with long spreading sequences,” *Journal of VLSI Signal Processing*, vol. 30, no. 1–3, pp. 273–291, 2002.
- [5] G. J. Foschini and M. J. Gans, “On limits of wireless communications in a fading environment when using multiple antennas,” *Wireless Personal Communications*, vol. 6, no. 3, pp. 311–335, 1998.
- [6] G. G. Raleigh and J. M. Cioffi, “Spatio-temporal coding for wireless communications,” *IEEE Trans. Communications*, vol. 46, no. 3, pp. 357–366, 1998.
- [7] D. Gesbert, H. Bolcskei, D. Gore, and A. Paulraj, “MIMO wireless channels: capacity and performance prediction,” in *IEEE Global Telecommunications Conference, 2000*, vol. 2, pp. 1083–1088, San Francisco, Calif, USA, December 2000.
- [8] A. Paulraj and T. Kailath, “Increasing capacity in wireless broadcast systems using distributed transmission/directional reception (DTDR),” U.S. Patent 5345599, Stanford University, Stanford, Calif, USA, September 1994.
- [9] G. J. Foschini, “Layered space-time architecture for wireless communication in a fading environment when using multiple antennas,” *Bell Labs Technical Journal*, vol. 1, no. 2, pp. 41–59, 1996.
- [10] V. Tarokh, N. Seshadri, and A. R. Calderbank, “Space-time codes for high data rate wireless communication: performance criterion and code construction,” *IEEE Transactions on Information Theory*, vol. 44, no. 2, pp. 744–765, 1998.
- [11] S. M. Alamouti, “A simple transmit diversity technique for wireless communications,” *IEEE Journal on Selected Areas in Communications*, vol. 16, no. 8, pp. 1451–1458, 1998.
- [12] V. Tarokh, H. Jafarkhani, and A. R. Calderbank, “Space-time block codes from orthogonal designs,” *IEEE Transactions on Information Theory*, vol. 45, no. 5, pp. 1456–1467, 1999.
- [13] Lindsog E. and A. Paulraj, “A transmit diversity scheme for channels with intersymbol interference,” in *Proc. IEEE Conference on International Communications*, vol. 1, pp. 307–311, New Orleans, La, USA, June 2000.
- [14] N. Al-Dhahir, “Single-carrier frequency-domain equalization for space-time block-coded transmissions over frequency-selective fading channels,” *IEEE Communications Letters*, vol. 5, no. 7, pp. 304–306, 2001.

- [15] S. Zhou and G. B. Giannakis, "Space-time coding with maximum diversity gains over frequency-selective fading channels," *IEEE Signal Processing Letters*, vol. 8, no. 10, pp. 269–272, 2001.
- [16] B. Hochwald, T. L. Marzetta, and C. B. Papadias, "A transmitter diversity scheme for wideband CDMA systems based on space-time spreading," *IEEE Journal on Selected Areas in Communications*, vol. 19, no. 1, pp. 48–60, 2001.
- [17] S. Barbarossa, G. Scutari, and A. Swami, "MUI-free CDMA systems incorporating space-time coding and channel shortening," in *Proc. IEEE Int. Conf. Acoustics, Speech, Signal Processing*, pp. 2213–2216, Orlando, Fla, USA, May 2002.
- [18] S. Zhou and G. B. Giannakis, "Single-carrier space-time block coded transmissions over frequency-selective fading channels," *IEEE Transactions on Information Theory*, vol. 49, no. 1, pp. 164–179, 2003.
- [19] H. Holma and A. Toskala, Eds., *WCDMA for UMTS: Radio Access for Third Generation Mobile Communications*, John Wiley & Sons, New York, NY, USA, 2001.
- [20] F. Petré, G. Leus, M. Engels, M. Moonen, and H. De Man, "Space-time chip equalization for WCDMA forward link with code-multiplexed pilot and soft handover," in *IEEE Global Telecommunications Conference, 2001*, vol. 1, pp. 280–284, San Antonio, Tex, USA, November 2001.
- [21] F. Petré, G. Leus, L. Deneire, M. Engels, and M. Moonen, "Adaptive space-time chip-level equalization for WCDMA downlink with code-multiplexed pilot and soft handover," in *Proc. IEEE Conference on International Communications*, pp. 1635–1639, New York City, NY, USA, April 2002.
- [22] J. G. Proakis, *Digital Communications*, McGraw-Hill, New York, NY, USA, 3rd edition, 1995.
- [23] M. Lenardi, A. Medles, and D. T. M. Slock, "Comparison of downlink transmit diversity schemes for RAKE and SINR maximizing receivers," in *Proc. IEEE Conference on International Communications*, pp. 1679–1683, New Orleans, La, USA, June 2001.
- [24] C. D. Frank, "MMSE reception of DS-SS-CDMA with open-loop transmit diversity," in *Proc. 2nd International Conference on 3G Mobile Communication Technologies, 2001*, pp. 156–160, London, UK, March 2001.

**Geert Leus** was born in Leuven, Belgium, in 1973. He received his Electrical Engineering degree and the Ph.D. degree in applied sciences from the Katholieke Universiteit Leuven, Belgium, in June 1996 and May 2000, respectively. He has been a Research Assistant and a Postdoctoral Fellow of the Fund for Scientific Research - Flanders, Belgium, from October 1996 till September 2003. During that period, Geert Leus was affiliated with the Electrical Engineering Department, the Katholieke Universiteit Leuven, Belgium. Currently, Geert Leus is an Assistant Professor at the Faculty of Electrical Engineering, Mathematics, and Computer Science, Delft University of Technology, The Netherlands. During the summer of 1998, he visited Stanford University, and from March 2001 till May 2002, he was a Visiting Researcher and Lecturer at the University of Minnesota. His research interests are in the area of signal processing for communications. Geert Leus received a 2002 IEEE Signal Processing Society Young Author Best Paper Award. He is a member of the IEEE Signal Processing for Communications Technical Committee, and an Associate Editor for the IEEE Transactions on Wireless Communications and the IEEE Signal Processing Letters.



**Frederik Petré** was born in Tienen, Belgium, 1974. He received the Electrical Engineering degree and the Ph.D. in Applied Sciences from the Katholieke Universiteit Leuven, Leuven, Belgium, in July 1997 and December 2003, respectively. In September 1997, he joined the Design Technology for Integrated Information and Communication Systems (DESICS) division, Interuniversity Micro-Electronics Center (IMEC), Leuven, Belgium. Within the Digital Broadband Terminals (DBATE) group of DESICS, he first performed predoctoral research on wireline transceiver design for twisted pair, coaxial cable, and powerline communications. During the fall of 1998, he visited the Information Systems Laboratory (ISL), Stanford University, California, USA, working on OFDM-based powerline communications. In January 1999, he joined the Wireless Systems (WISE) group of DESICS as a Ph.D. Researcher, funded by the Institute for Scientific and Technological Research in Flanders (IWT). Since January 2004, he is a Senior Scientist within the Wireless Research group of DESICS. He is investigating the baseband signal processing algorithms and architectures for future wireless communication systems, like third generation (3G) and fourth generation (4G) cellular networks, and wireless local area networks (WLANs). His main research interests are modulation theory, multiple access schemes, channel estimation and equalization, smart antenna, and MIMO techniques. He is a member of the ProRISC technical program committee and the IEEE Benelux Section on Communications and Vehicular Technology (CVT). He is a member of the Executive Board and Project Leader of the Reconfigurable Radio project of the Network of Excellence in Wireless Communications (NEWCOM), established under the sixth framework of the European Commission.



**Marc Moonen** received the Electrical Engineering degree and the Ph.D. degree in applied sciences from the Katholieke Universiteit Leuven, Leuven, Belgium, in 1986 and 1990, respectively. Since 2000, he has been an Associate Professor at the Electrical Engineering Department, Katholieke Universiteit Leuven, where he is currently heading a research team of sixteen Ph.D. candidates and postdocs working in the area of signal processing for digital communications, wireless communications, DSL, and audio signal processing. He received the 1994 KU Leuven Research Council Award, the 1997 Alcatel Bell (Belgium) Award (with Piet Vandaele), and was a 1997 "Laureate of the Belgium Royal Academy of Science." He was the Chairman of the IEEE Benelux Signal Processing Chapter (1998–2002), and is currently a EURASIP AdCom Member (European Association for Signal, Speech, and Image Processing, 2000). He is the Editor-in-Chief for the EURASIP Journal on Applied Signal Processing (2003), and a member of the editorial board of Integration, the VLSI Journal, IEEE Transactions on Circuits and Systems II, and IEEE Signal Processing Magazine.

